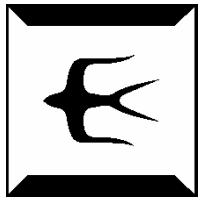




Safety Analysis of Vehicle Platoon under V2V2I Communication



Takuto Takagi

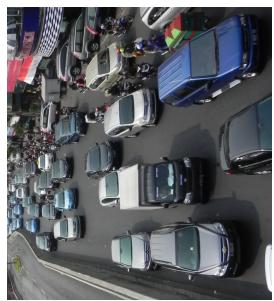
FL11-15-1

20th, December, 2011



Introduction

Background



- Traffic congestion[1]
- Adaptive cruise control[2]
- Intelligent road transportation system[8]
 - Vehicle-to-Vehicle communication(V2V)
 - Vehicle-to-Infrastructure communication(V2I)

Approaches

- Macro perspective : On-ramp control[3,4], Transportation network[5]
- Micro perspective : Vehicle platoon control[6,7]



Purpose of Research

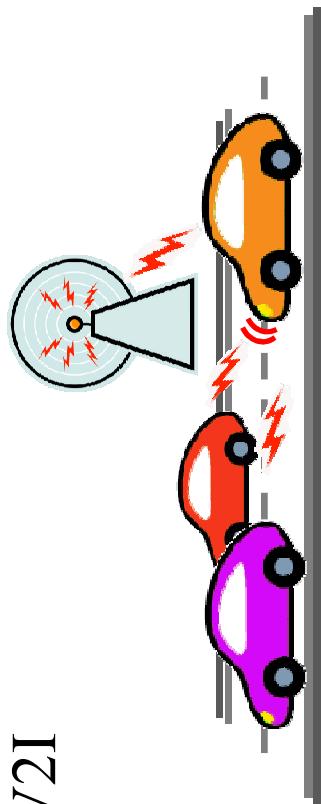
New approach

Middle perspective

- Considering **vehicular strings** and **infrastructure**
- Under Vehicle-to-Vehicle-to-Infrastructure communication

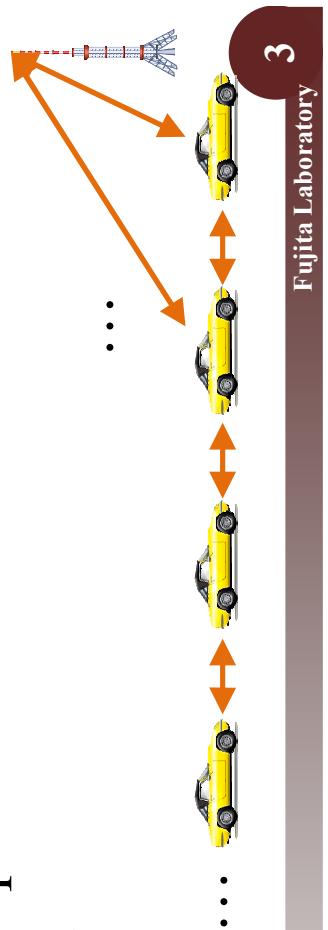
Vehicle-to-Vehicle-to-Infrastructure(V2V2I) communication

- The **hybrid system** of the V2V and V2I
- Vehicles are controlled by
“Local” and **“Global”** information



Objectives

- Problem description of middle perspective
- Safety analysis of vehicle platoon
- Simulation





Problem Description

Vehicle platoon model[3,4]

$$\begin{bmatrix} x(k+1) \\ v(k+1) \end{bmatrix} = \begin{bmatrix} I & T \\ 0 & I \end{bmatrix} \begin{bmatrix} x(k) \\ v(k) \end{bmatrix} + \begin{bmatrix} 0 \\ I \end{bmatrix} u(k)$$

$$x(k) = [x_1(k), \dots, x_n(k)]^T \quad u(k) = [u_1(k), \dots, u_n(k)]^T$$

$$v(k) = [v_1(k), \dots, v_n(k)]^T \quad T : \text{Sampling time}$$

Error[9,10]

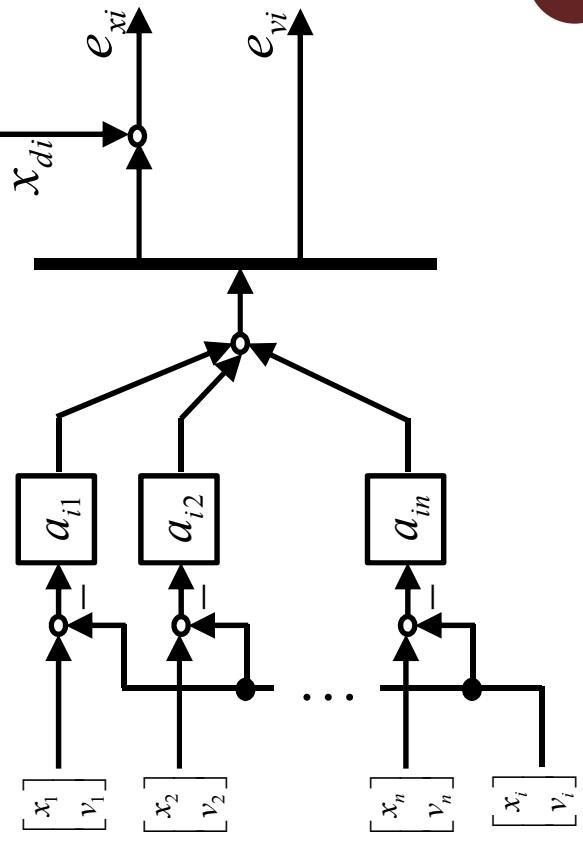
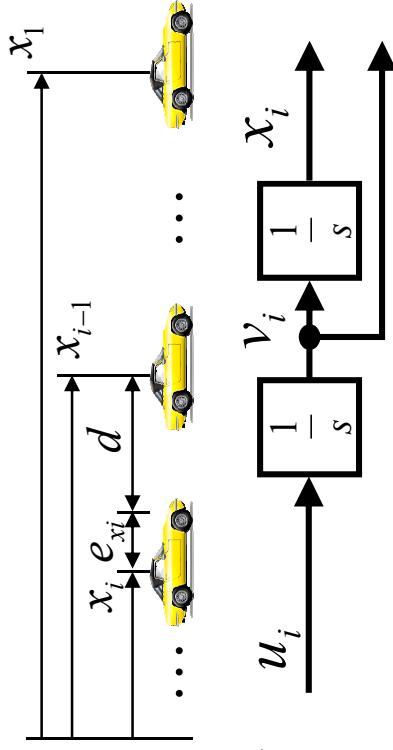
*i*th vehicle's position and velocity error

$$\begin{cases} e_{xi}(k) = \sum_{j=1}^n \alpha_{ij} (x_j(k) - x_i(k)) + x_{di} \\ e_{vi}(k) = \sum_{j=1}^n \alpha_{ij} (v_j(k) - v_i(k)) \end{cases}$$

x_{di} : Constant desired spacing

Weighted communication state

$$\begin{cases} 0 \leq \alpha_{ij} \leq 1 \\ \alpha_{ii} = -1 \end{cases} \quad \sum_{j \neq i} \alpha_{ij} = 1 \quad \sum_{j=1}^n \alpha_{ij} = 0$$





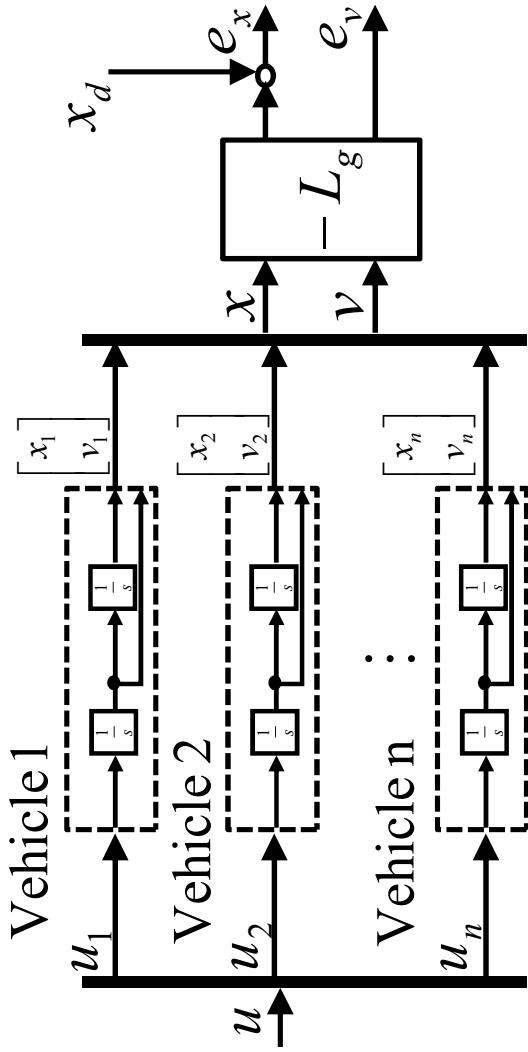
Problem Description

Platoon's position and velocity error

$$\begin{cases} e_x(k) = -L_g x(k) + x_d \\ e_v(k) = -L_g v(k) \end{cases}$$

$$x_d(k) = [x_{d1}(k), \dots, x_{dn}(k)]^T$$

$$-L_g = \begin{bmatrix} -1 & \alpha_{12} & \cdots & \alpha_{1n} \\ \alpha_{21} & -1 & \cdots & \alpha_{2n} \\ \vdots & & & \\ \alpha_{n1} & \alpha_{n2} & \cdots & -1 \end{bmatrix}$$



Input from Infrastructure

ith vehicle's input

$$u_i(k) = k'_{\nu} e_{xi}(k) + k'_{\nu} e_{vi}(k)$$

$$k'_{\nu} = k_{\nu} T \quad k'_{\mu} = k_{\mu} T$$

Reference relative position

ith vehicle's input

$$x_{di} = \sum_{j=1}^n j a_{ij} d(k) \quad \mathcal{I} = [1, 2, 3, \dots]^T$$

$$d(k) \quad d(k) \quad d(k)$$



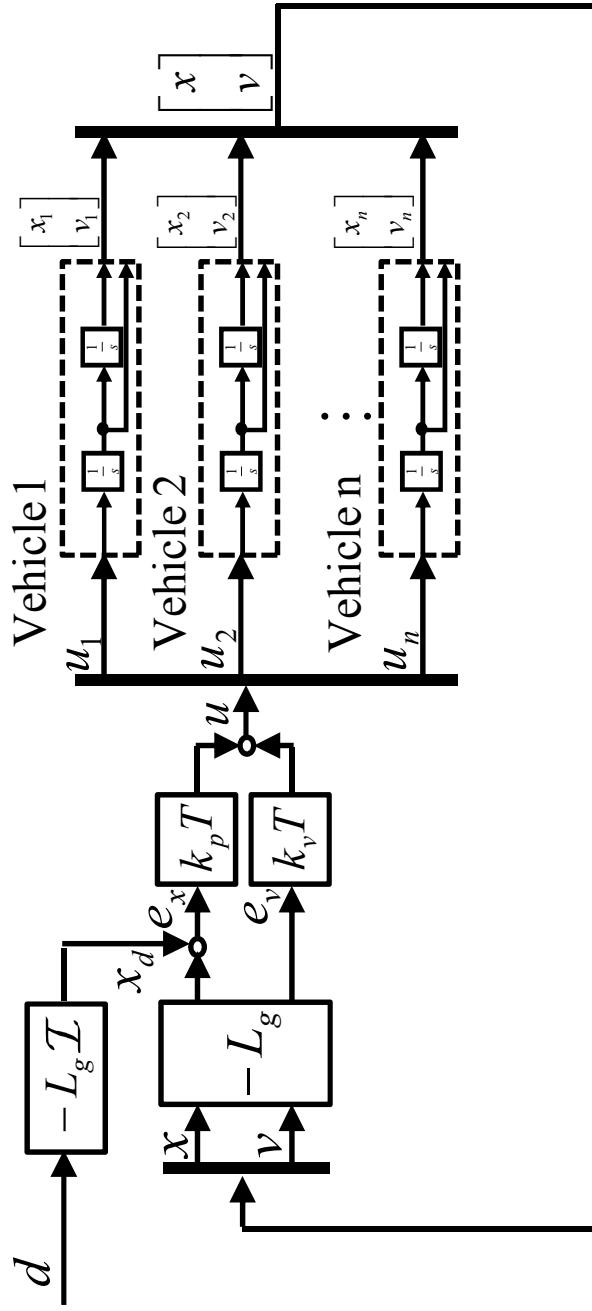
Platoon's input

$$u(k) = k'_{\mu} e_x(k) + k'_{\nu} e_{\nu}(k)$$

Platoon's input

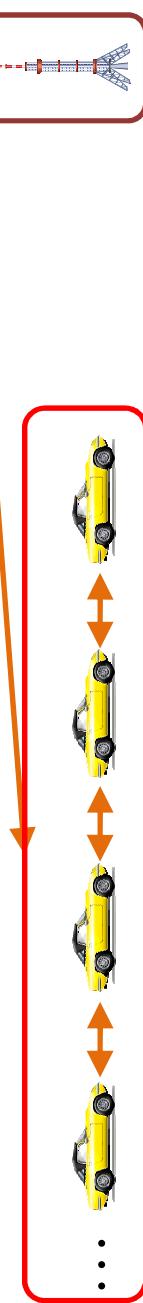


Problem Description



State equation of vehicle platoon

$$\begin{bmatrix} x(k+1) \\ v(k+1) \end{bmatrix} = \begin{bmatrix} I & TI \\ -k'_p L_g & I - k'_v L_g \end{bmatrix} \begin{bmatrix} x(k) \\ v(k) \end{bmatrix} - k'_p \begin{bmatrix} 0 \\ L_g \mathcal{I} \end{bmatrix} d(k) \quad \dots(1)$$



Vehicle platoon

Input from Infrastructure



Safety Analysis

Definition of safety

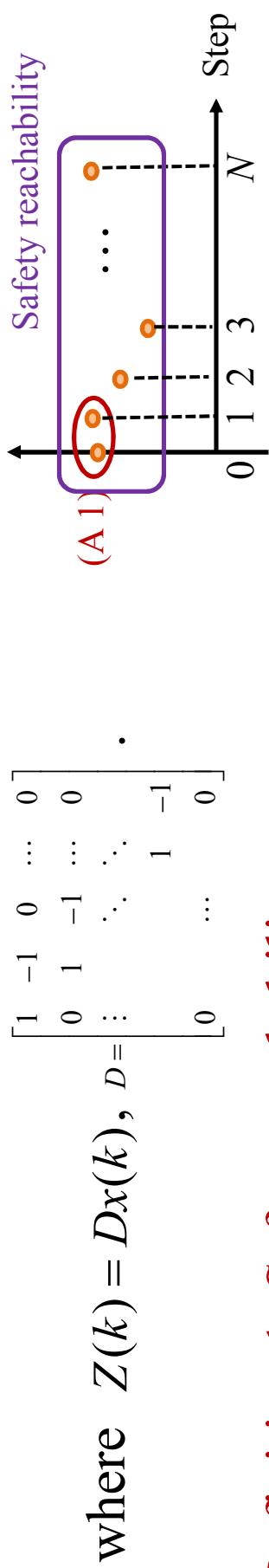
$$\begin{bmatrix} x(k+1) \\ v(k+1) \end{bmatrix} = \begin{bmatrix} x(k) + T v(k) \\ v(k) - k'_{\rho} L_g x(k) - k'_{\rho} L_g v(k) - k'_{\rho} L_g \mathcal{Z} d(k) \end{bmatrix}$$

$$\begin{bmatrix} x(k+2) \\ v(k+2) \end{bmatrix} = \begin{bmatrix} x(k) + 2 T v(k) - T(k'_{\rho} L_g x(k) + k'_{\nu} L_g v(k)) - k'_{\rho} T L_g \mathcal{Z} d(k) \\ -k'_{\rho} L_g (x(k) + T v(k)) + k'_{\nu} L_g (k'_{\rho} L_g x(k) + k'_{\nu} L_g v(k) + k'_{\rho} L_g \mathcal{Z} d(k)) - k'_{\rho} T L_g \mathcal{Z} d(k) \end{bmatrix}$$

Assumption (A1):

The system (1) satisfies $Z(0) > 0, Z(1) > 0,$

ex) $n=2$



Definition 1: Safety reachability

The system (1) is **safety reachable** if given any $Z(0)$ satisfying (A1)

there exists a $d(k), \forall 0, 1, \dots, N-2$ such that

$$Z(k) > 0, \forall 2, 3, \dots, N.$$



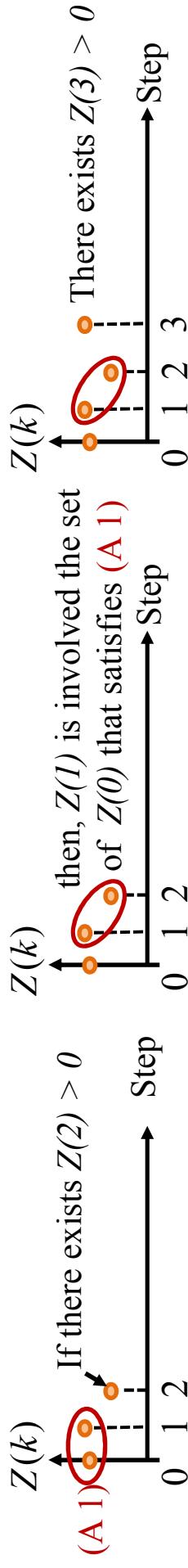
Safety Analysis

Main result

Theorem 1:

The system (1) is **safety reachable** iff given any $Z(0)$ satisfies (A1)
there exists a $d(0) > 0$ that satisfies $Z(2) > 0$

Sketch of proof: Sufficient condition



Theorem 2:

In order for given any $Z(0)$ satisfies (A1) there exists a $d(0) > 0$ that satisfies $Z(2) > 0$, it is necessary that any $i = 1, 2, \dots, n-1$ satisfies

$$\{-DL_g\mathcal{Z}\}_i \neq 0.$$

Sketch of proof:

$Z(2) = Dx(2) = Dx(0) + 2T Dv(0) - T(k'_{\rho} DL_g x(0) + k'_{\nu} DL_g v(0)) - k'_{\rho} TDL_g \mathcal{Z} d(0)$
If there exists i satisfying $\{-DL_g\mathcal{Z}\}_i = 0, x_i(2) - x_{i+1}(2)$ doesn't depend on $d(0)$



Safety Analysis

Safety unreachability communication structure

Considering the case $\{-DL_g\mathcal{Z}\}_i = 0$

ex) Predecessor following

$$-DL_g\mathcal{Z} = D \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & -1 & & \\ & 1 & -1 & \\ & & 1 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 2 \\ 3 \\ 4 \end{bmatrix} = D \begin{bmatrix} 0 \\ -1 \\ -1 \\ -1 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$



Corollary 1:

The system (1) is safety unreachable iff given $i = 1, 2, \dots, n-1$

$-L_g$ satisfies the following condition

$$\begin{bmatrix} a_{i(i-1)} - a_{i(i+1)} & a_{i(i-2)} - a_{i(i+2)} & \cdots \end{bmatrix} = \begin{bmatrix} a_{(i+1)i} - a_{(i+1)(i+2)} & a_{(i+1)(i-1)} - a_{(i+1)(i+3)} & \cdots \end{bmatrix}$$

with $a_{ij} = 0, j < 0$.

Sketch of proof:

$$\begin{aligned} \{-L_g\mathcal{Z}\}_i &= \sum_{j=1}^n ja_{ij} - i \sum_{j=1}^n a_{ij} = -(a_{i(i-1)} - a_{i(i+1)}) - 2(a_{i(i-2)} - a_{i(i+2)}) \dots \\ \{-DL_g\mathcal{Z}\}_i &= \sum_{j=1}^n j(a_{ij} - a_{(i+1)j}) - \sum_{j=1}^n ia_{ij} - (i+1)a_{(i+1)j} = 0 \end{aligned}$$



Relatively-equal communication structure makes the system (1) safety unreachable



Safety Analysis

Safety reachability communication structure

There exists communication structures satisfying $\{-DL_g \mathcal{Z}\}_i \neq 0$

ex) Leader and Predecessor following

$$-DL_g \mathcal{Z} = D \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 1 & -1 & & & \\ 1/2 & 1/2 & -1 & & \\ 1/2 & 1/2 & -1 & & \\ 1/2 & 1/2 & -1 & & \end{bmatrix} \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{bmatrix} = D \begin{bmatrix} 0 \\ -1 \\ -3/2 \\ -2 \\ -5/2 \end{bmatrix} = \begin{bmatrix} 1 \\ 1/2 \\ 1/2 \\ 1/2 \\ 0 \end{bmatrix}$$

Theorem 3:

If $-L_g$ satisfies the following condition:

$$\sum_{j=1}^n j a_{1j} > \dots > \sum_{j=1}^n j a_{nj}$$

Then there exists a $d(0) > 0$ that satisfies $Z(2) > 0$.

Sketch of proof:

$$\begin{aligned} Z(2) &= Dx(0) + 2TDv(0) - T(k'_p DL_g x(0) + k'_v DL_g v(0)) - k'_p TDL_g \mathcal{Z} d(0) > 0 \\ \Updownarrow -DL_g \mathcal{Z} d(0) &\geq -\frac{D}{k'_p T} (x(0) + 2Tv(0) - k'_p TL_g x(0) - k'_v TL_g v(0)) \end{aligned}$$

If given any $i = 1, 2, \dots, n-1$ satisfies $\{-DL_g \mathcal{Z}\}_i > 0$, $d(0)$ has an only lower bound

Latter vehicle should lay more weight on the leader information





Simulation

Comparative Verification

Safety unreachability ex:

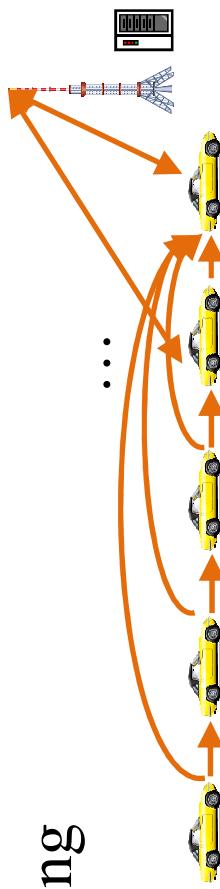
Predecessor following



Safety reachability ex:

Leader and Predecessor following

- Constant control
- Infrastructure control



Simulation Settings

Constant control

$$d(k) = d_{ref}$$

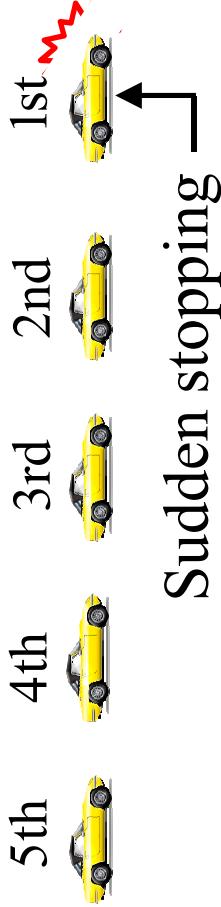
Infrastructure control

$$d(k) = \begin{cases} d_{ref} & (d_{ref} \geq d_{min}(k)) \\ d_{min}(k) & (d_{min}(k) > d_{ref}) \end{cases}$$

$$d_{min}(k) = \min_{d(k)} d(k) + \delta$$

$$\text{s.t. } -(d(k) + \delta) D L_g \mathcal{Z} \geq -\frac{D}{k_p T} (x(k) + 2T v(k) - k_p^T T L_g x(k) - k_p^T T L_g v(k))$$

Speed profile



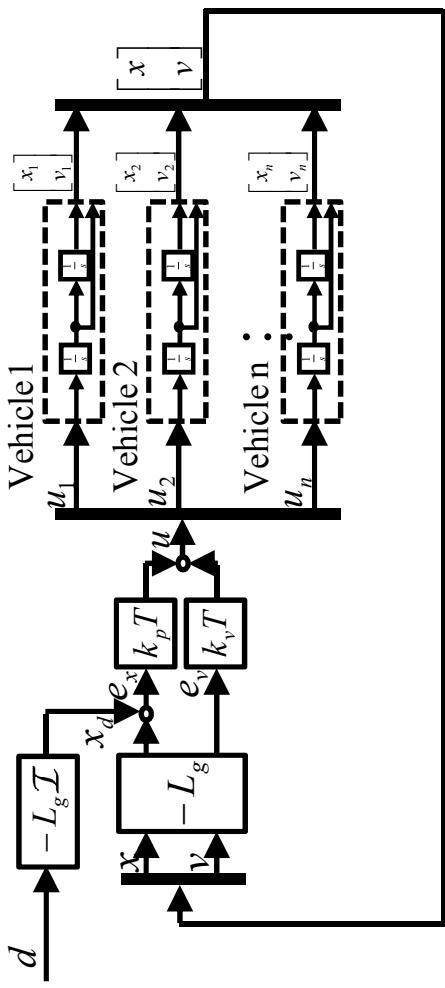
Sudden stopping



Summary

Summary

- Problem description of middle perspective



$$\begin{bmatrix} x(k+1) \\ v(k+1) \end{bmatrix} = \begin{bmatrix} I & TI \\ -k_v^T L_g & I - k_v^T L_g \end{bmatrix} \begin{bmatrix} x(k) \\ v(k) \end{bmatrix} - k_v^P \begin{bmatrix} 0 \\ L_g Z \end{bmatrix} d(k)$$

- Safety analysis of vehicle platoon

Theorem 1

Theorem 2

Corollary 1

Theorem 3



Reference

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Appendix



Simulation

Comparative Verification

Precede following

$$L_g = \begin{bmatrix} \ddots & & & \\ & 1 & -1 & \ddots \end{bmatrix}$$

$$L_g = \begin{bmatrix} \vdots & \ddots & \ddots & \vdots \\ 1/2 & 0 & \dots & 1/2 & -1 & 0 \\ 1/2 & 0 & \dots & 0 & 1/2 & -1 \\ \vdots & & & & \ddots & \vdots \end{bmatrix}$$



Simulation Settings

Parameter

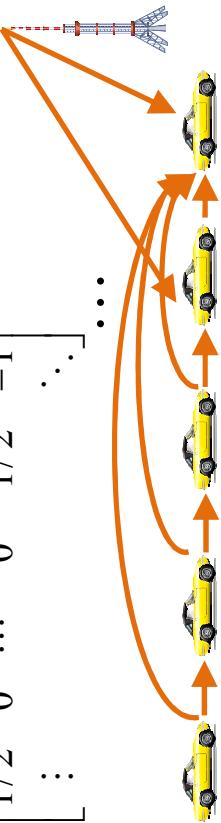
$$k_p = 2, k_v = 2$$

$d = 25[\text{m}] \rightarrow$ Time head way

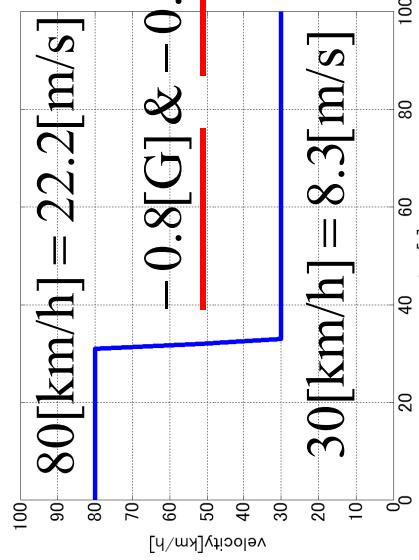
$$80[\text{km/h}]: 1.13[\text{s}]$$

$$30[\text{km/h}]: 3.00[\text{s}]$$

Leader and Precede following



Speed profile

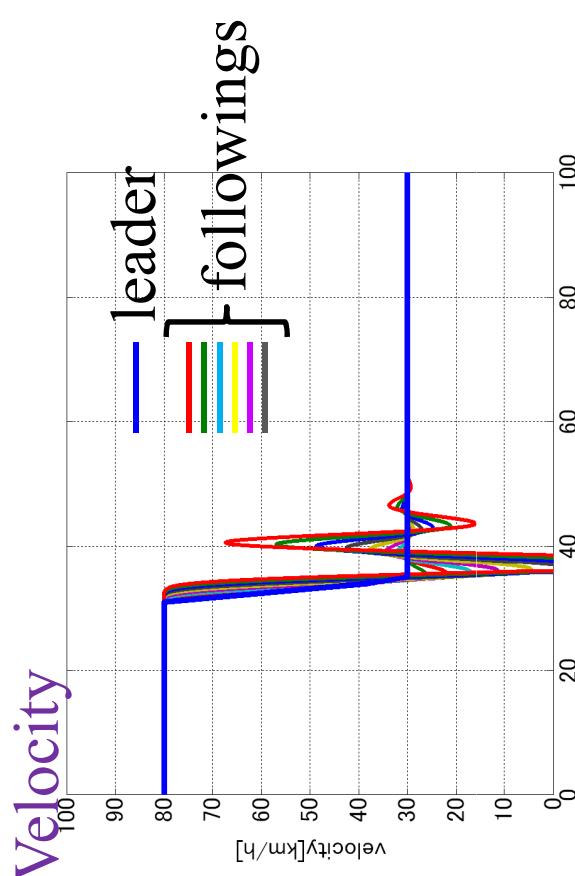




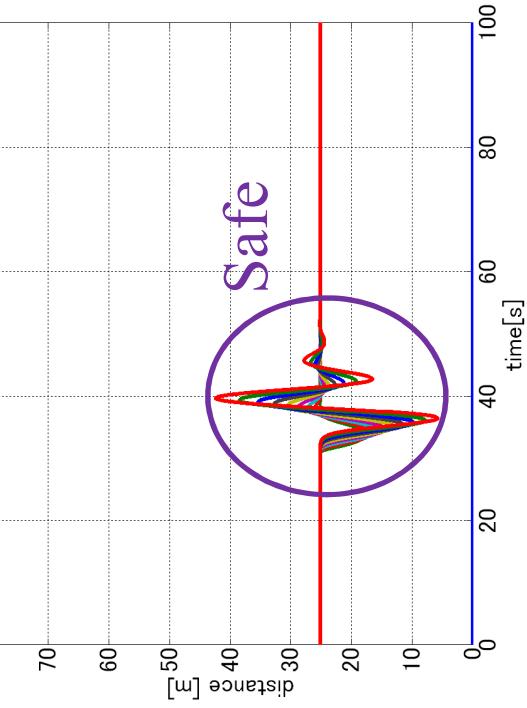
Simulation Result (-0.4[G])

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Precede following



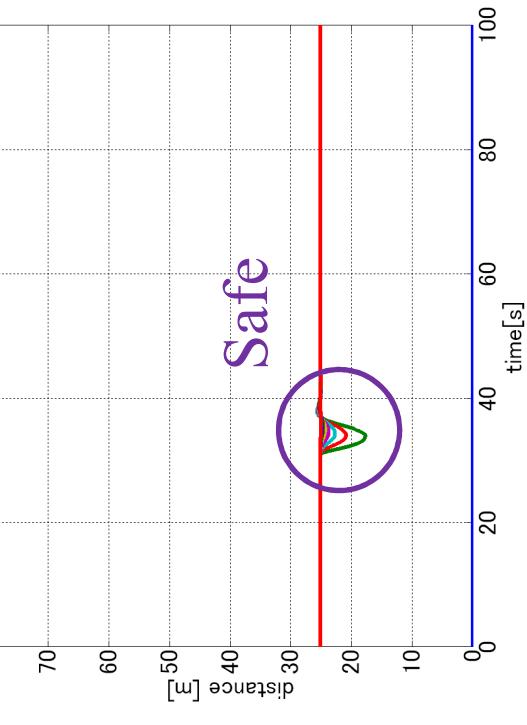
Relative distance



Leader and Precede following



Relative distance

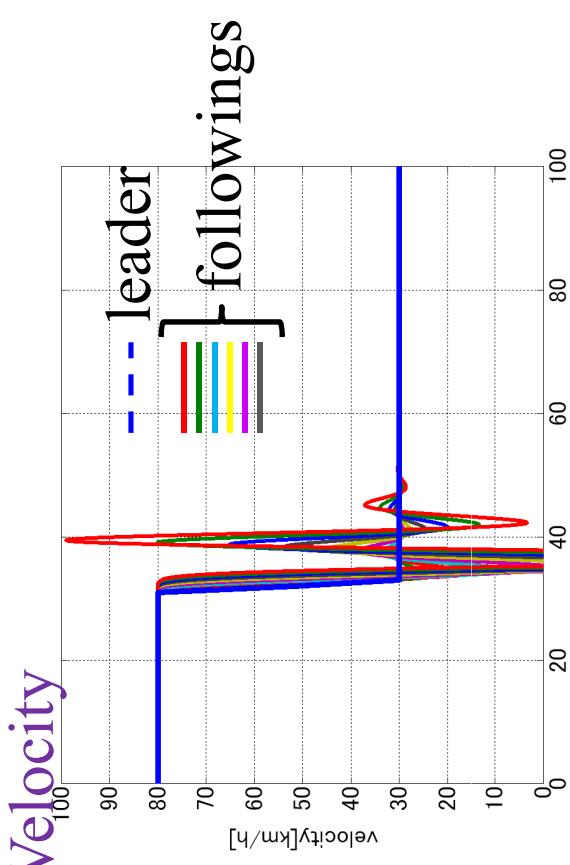




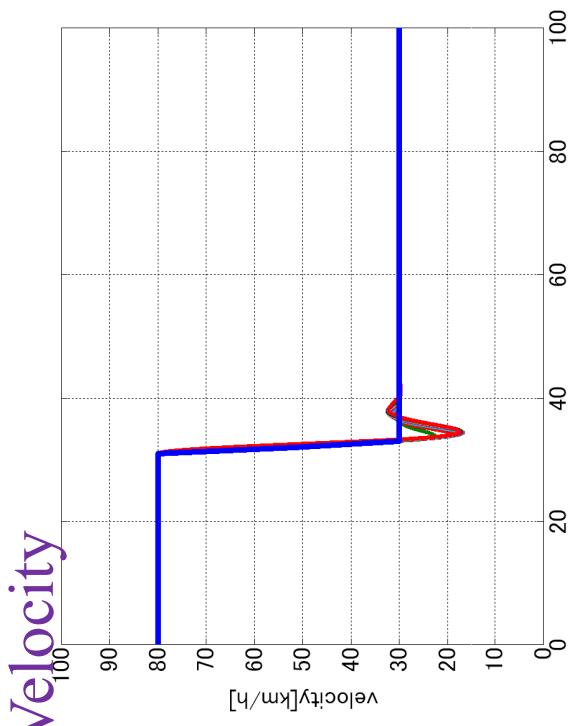
Simulation Result (-0.8[G])

Tokyo Institute of Technology

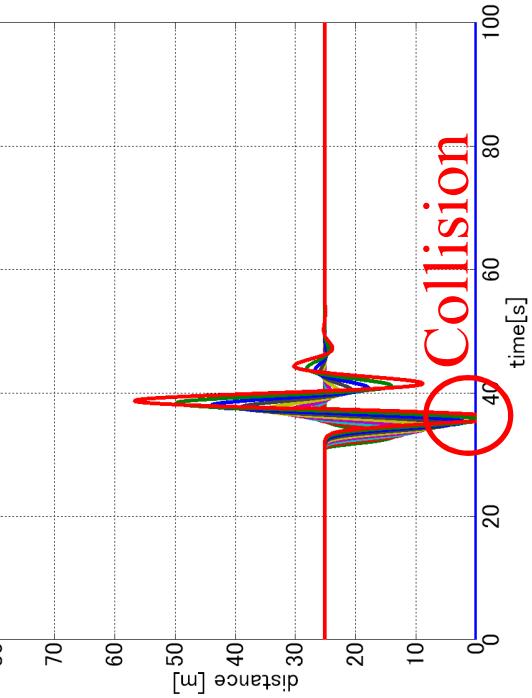
Precede following



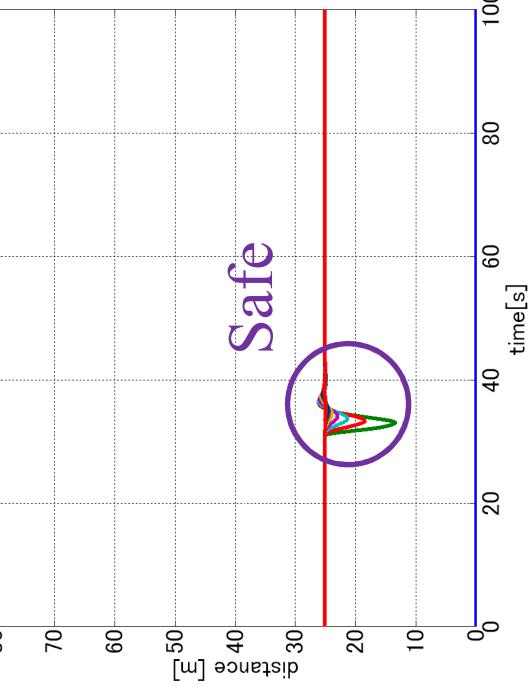
Leader and Precede following



Relative distance



Relative distance



Relative distance



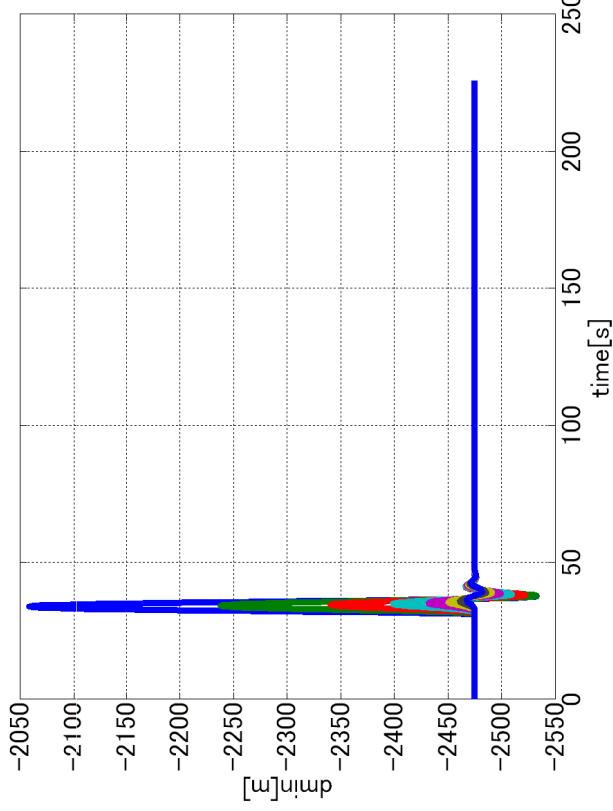
Simulation Result

Leader and Precede following

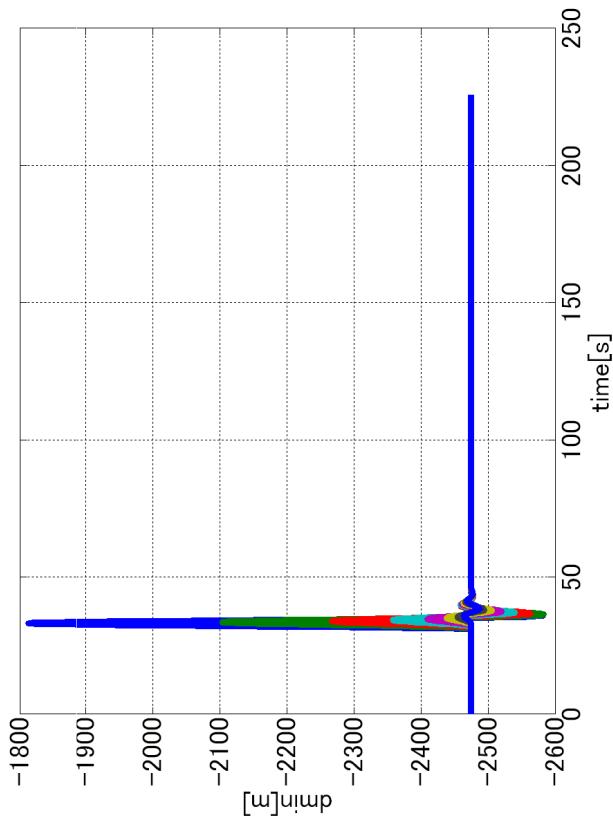
$$dDL_g J > -\frac{D}{k'_{_P} t} \left(x(k) + 2t\nu(k) + k'_{_P} t L_g x(k) + k'_{_V} t L_g \nu(k) \right)$$

$$\uparrow d_{\min} DL_g J = -\frac{D}{k'_{_P} t} \left(x(k) + 2t\nu(k) + k'_{_P} t L_g x(k) + k'_{_V} t L_g \nu(k) \right)$$

dmin -0.4[G]



dmin -0.8[G]



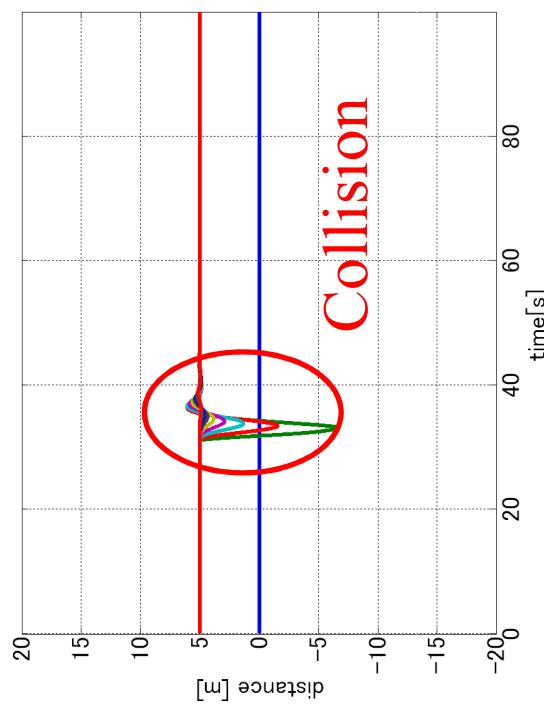
dmin has much luxury



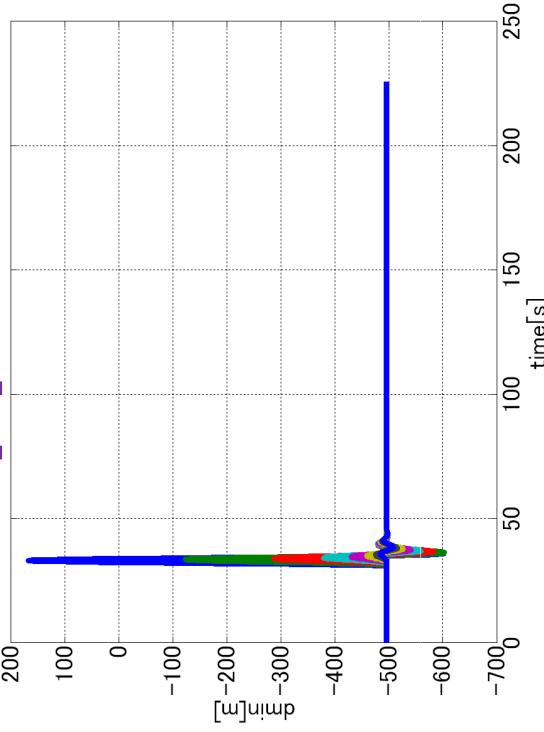
Simulation Result(-0.8[G])

$$d = 25[\text{m}] \rightarrow 5[\text{m}]$$

Relative distance



d_{min}-0.8[G]



Simulation Result(-0.8[G])

d_{min} feedback vehicle control

$$\begin{bmatrix} x(k+1) \\ v(k+1) \end{bmatrix} = \begin{bmatrix} I & tI \\ k'_{\nu} L_g & I + k'_{\nu} L_g \end{bmatrix} \begin{bmatrix} x(k) \\ v(k) \end{bmatrix} + k'_{\nu} \begin{bmatrix} 0 \\ L_g J \end{bmatrix} d'_{\min}$$

$$d'_{\min} = \begin{cases} d_{\min} & d_{\min} > d \\ d & d_{\min} < d \end{cases}$$



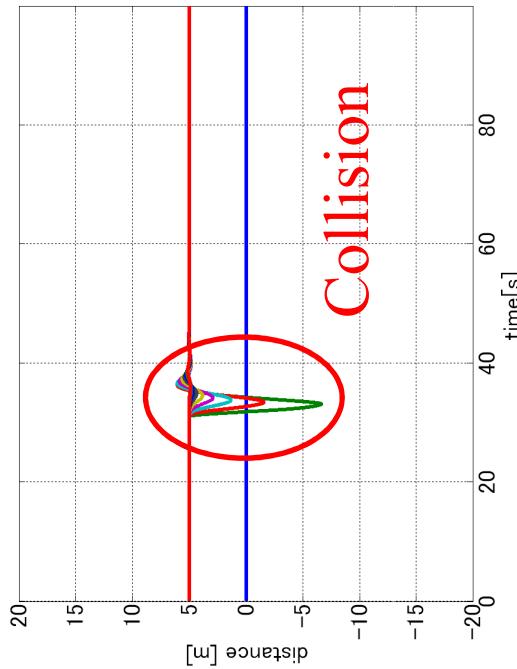
Simulation Result(-0.8[G]) $d = 5[m]$

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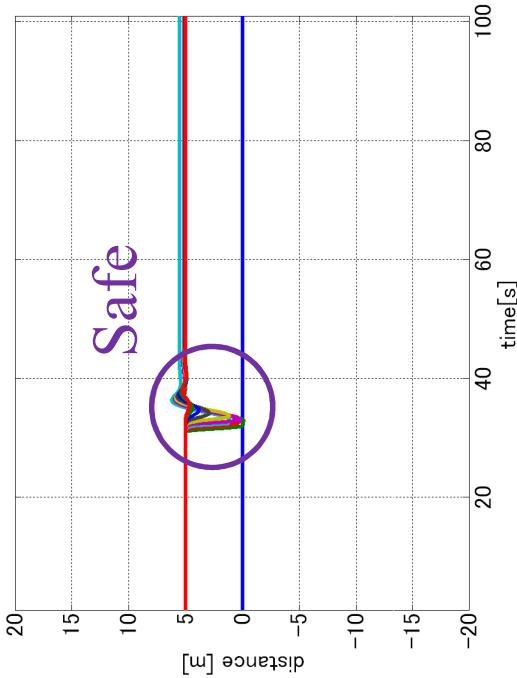
Constant d vehicle control

dmin feedback vehicle control

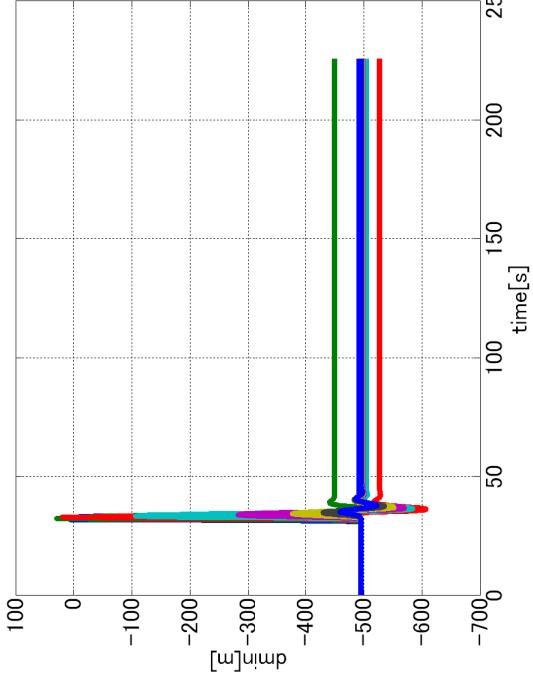
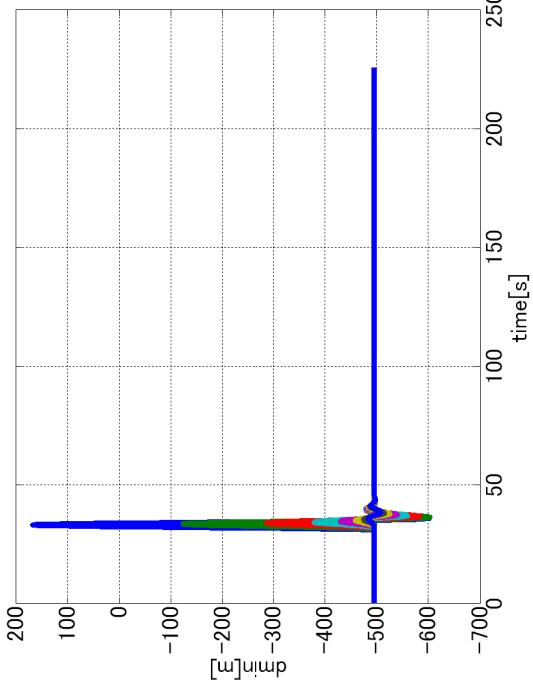
Relative distance



Relative distance



Toky



Toky