

Optimal Power Flow of Power Networks with Battery via Distributed Predictive Control



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Outline

- Introduction
- Power Network
- Predictive Control
- Distributed Control
- Simulation



Introduction

Power Network

a network to send from various power supplies to each consumer

Optimal Power Flow(OPF)

optimize a certain objective over power network variables

■ constraints

(bounds on voltages, line loading, etc.)

→ we must optimize **under certain constraints**

■ renewable energy

(solar, or wind power, etc.)

output fluctuates widely and randomly, but predictably

→ formulate an optimal power flow problem with **battery**



Introduction

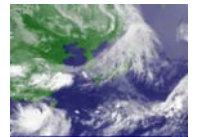
OPF has been studied *static optimization*, but...

Introduction of battery

→ **Dynamics of battery in OPF**

Desired Control

formulate a OPF of power network with storage across multiple time periods



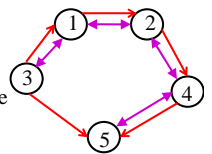
distributed predictive control

- applicable {
 - predictive information in the future
 - constraints

- only by local information

objective of this work

Optimize power consumption by applying the distributed predictive control to power network



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Power Network - example

Node

- Demand D_1, D_2, D_3
- Generation G_1
- Battery B_1, B_2, B_3
- Calculator
- Renewable energy R_1, R_2, R_3

Link

$$\mathcal{E}_{DG} = \{(G_1, D_1), (G_1, D_2), (G_1, D_3)\}$$

$$\mathcal{E}_{DR} = \{(R_1, D_1), (R_2, D_2), (R_3, D_3)\}$$

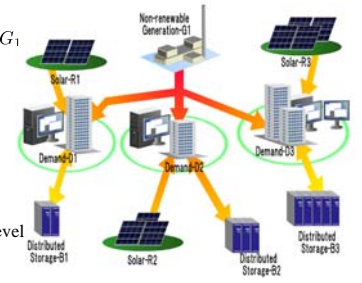
$$\mathcal{E}_{DB} = \{(B_1, D_1), (B_2, D_2), (B_3, D_3)\}$$

Utility

- move battery level closer to target level
- curb costs
- equate supply and demand

Constraints

- Cannot charge battery more than the supply by the others
- Supply is always nonnegative (without battery)
- about battery level
- Upper bound M in the sum of the supply to some demands





Power Network

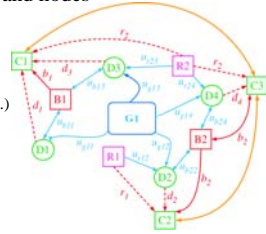
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Network structure

be composed of several different links and nodes

Node $\mathcal{V} = \mathcal{D} \cup \mathcal{G} \cup \mathcal{R} \cup \mathcal{B} \cup \mathcal{C}$

- Demand \mathcal{D} — customer
- Generation \mathcal{G} — nonrenewable generation (Gas, Coal, Nuclear, Hydro, etc.)
- Renewable energy \mathcal{R} — renewable, variable (Solar, Wind, tidal force, etc.)
- Battery \mathcal{B} — distributed storage
- Calculator \mathcal{C} — scheduling flows of electricity



Link $\mathcal{E} \subset \mathcal{V} \times \mathcal{V}$

- energy flow \mathcal{E}_E
 - to demand \mathcal{E}_D (blue arrow)
 - to battery \mathcal{E}_B (cyan arrow)
- information flow \mathcal{E}_I
 - battery info \mathcal{E}_S (red arrow)
 - Predictive info \mathcal{E}_P (dashed red arrow)
 - communication info \mathcal{E}_C (yellow arrow)

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7

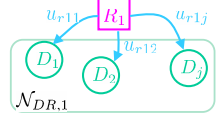


Power Network

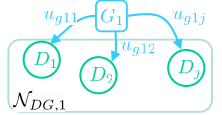
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Signal $k = 0, 1, 2, 3, \dots$

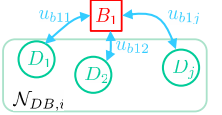
- Renewable energy $u_{r ij}, D_j \in \mathcal{N}_{DR,i}$
 $\mathcal{N}_{DR,i} = \{D_j \in \mathcal{D} \mid (R_i, D_j) \in \mathcal{E}_{DR}\}$
 $: D_j$ set that can be supplied by Renewable energy



- Generator $u_{g ij}, D_j \in \mathcal{N}_{DG,i}$
 $\mathcal{N}_{DG,i} = \{D_j \in \mathcal{D} \mid (G_i, D_j) \in \mathcal{E}_{GR}\}$
 $: D_j$ set that can be supplied by Generator



- Battery $u_{b ij}, D_j \in \mathcal{N}_{DB,i}$
 $(B_i, D_j) \in \mathcal{E}_{DB} \Leftrightarrow (D_j, B_i) \in \mathcal{E}_B$
Demand that can be supplied by battery can charge the battery oppositely



$$\mathcal{N}_{DB,i} = \{D_j \in \mathcal{D} \mid (B_i, D_j) \in \mathcal{E}_{DB}\}$$

$: D_j$ set that can be supplied by battery

$$\mathcal{N}_{BD,i} = \{B_i \in \mathcal{B} \mid (D_j, B_i) \in \mathcal{E}_B\}$$

$: B_i$ set that can be supplied by Generator

$u_{b ij} < 0$ means battery charge

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8

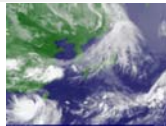


Power Network

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Problem of this system

- Instability of renewable energy
Mainly controlled by the weather
It is possible to predict at some level



Dynamics of battery

$$\Sigma_i : b_i(k+1) = b_i(k) - \sum_{j \in \mathcal{N}_{DB,i}} u_{bij}, i \in \mathcal{B}, \dots (1)$$

Past information

$b_i(k)$: remaining battery level

There is relativity around the time of a time period

We must use dynamic optimization across multiple time periods

Predictive control

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9



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- Introduction
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10



predictive control

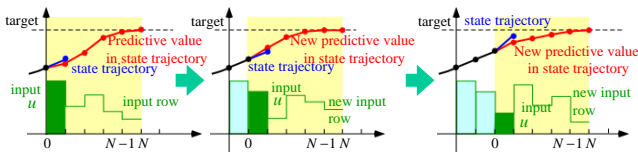
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Utility function $\min_{U^N} J = (U^N)^T H U^N + h(b_0) U^N$

Constraints $s.t. AU^N \leq B(b_0)$

Predictive control

- solve optimal control of finite horizon by the state x_0 at the time k
- input the first step of the obtained input row
- in the next time step, return 1) and repeat



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11



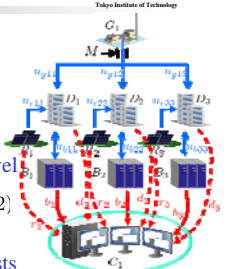
predictive control

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Utility function

utility requested on power network

$$\min_{u(k), k=0, \dots, N} J = q_d J_d + q_b J_b + q_g J_g$$



Bring remaining battery level close to target level

$$J_b = \sum_{k=0}^N \sum_{i \in \mathcal{B}} w_{B,i} (b_i(k) - b_{i,ref})^2 \dots (2)$$

battery level target level

Curb power feeding from generation for the costs

$$J_g = \sum_{k=0}^N \sum_{i \in \mathcal{G}} \sum_{j \in \mathcal{N}_{DG,i}} w_{G,ij} u_{gij}^2(k) \dots (3)$$

Bring supply close to demand

$$J_d = \sum_{k=0}^N \sum_{j \in \mathcal{D}} w_{D,j} \left(\sum_{i \in \mathcal{N}_{GD,j}} u_{gij}(k) + \sum_{i \in \mathcal{N}_{RD,j}} u_{rij}(k) + \sum_{i \in \mathcal{N}_{BD,j}} u_{bij}(k) - d_j(k) \right)^2 \dots (4)$$

Generator Renewable energy Battery Demand

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12



predictive control

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Constraints

- Supply by generator and renewable energy is always nonnegative

$$u_{gij}(k) \geq 0, u_{rij}(k) \geq 0 \quad \dots (5)$$

- Cannot charge battery more than the supply by the others

$$\sum_{i \in \mathcal{N}_{BD,j}} u_{bij}(k) \geq -\sum_{i \in \mathcal{N}_{GD,j}} u_{gij}(k) - \sum_{i \in \mathcal{N}_{RD,j}} u_{rij}(k) \quad \dots (6)$$

- Renewable energy is distributed by demands in $\mathcal{N}_{DR,i}$

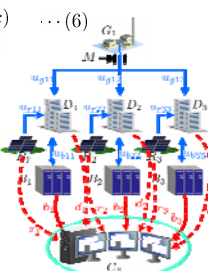
$$r_i(k) = \sum_{j \in \mathcal{N}_{DR,i}} u_{rij}(k) \quad \dots (7)$$

- Upper bound in the sum of the supply to some demands

$$\sum_{j \in \mathcal{D}_{sum,l}} u_{gij} \leq M, M \geq 0 \quad \dots (8)$$

- Maximum battery level

$$b_i(k) \in [0, B_{i,max}] \quad \dots (9)$$



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13



predictive control

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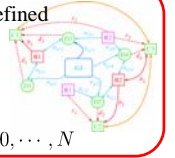
Predictive control problem

All the utility functions and the constraints are defined by convex function, so the combination is free

$$\min_{u(k), k=0, \dots, N} J = q_d J_d + q_b J_b + q_g J_g$$

subject to (1), (5) - (9).

$$b(0) = b_0, r_i(k) = r_{i,k}, d_j(k) = d_{j,k}, i \in \mathcal{R}, j \in \mathcal{D}, k = 0, \dots, N$$



Quadratic programming problem

Utility function $\min_{U^N} J = (U^N)^T H U^N + h(b_0) U^N$
 Constraints $A U^N \leq B(b_0)$

optimal control of finite horizon
 method : Primal-Dual Algorithm

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14



predictive control

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Primal-Dual Algorithm

Primal : $\min_{U^N} J = (U^N)^T H U^N + h(b_0) U^N \quad A U^N \leq B(b_0)$

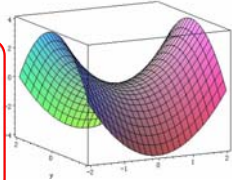
Lagrange-Dual $L(U^N, \lambda) = (U^N)^T H U^N + h(b_0) U^N - \lambda(A U^N - B(b_0))$

Dual : $\max_{\lambda > 0} D(\lambda) = \max_{\lambda > 0} \{ \min_{U^N} L(U^N, \lambda) \}$

Primal : convex function
 Dual : concave function

Optimized solution is the saddle point by both variables

$$\hat{U}^N = \left(\frac{\partial D}{\partial U^N} \right)_{U^N, \lambda} \quad \text{Update alternately} \quad \hat{\lambda} = - \left(\frac{\partial D}{\partial \lambda} \right)_{U^N, \lambda}$$



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15



Outline

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- Introduction
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16



Distributed Control

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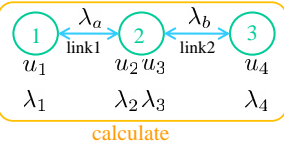
Power network tends to become large-scale

→ The amount of calculation becomes huge

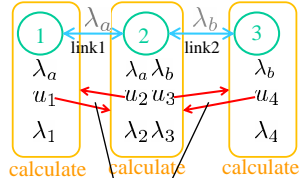
Distributed Control

Each grid controls only by local information

Centralized control



Distributed control



communicate relevant information to decide λ_a, λ_b

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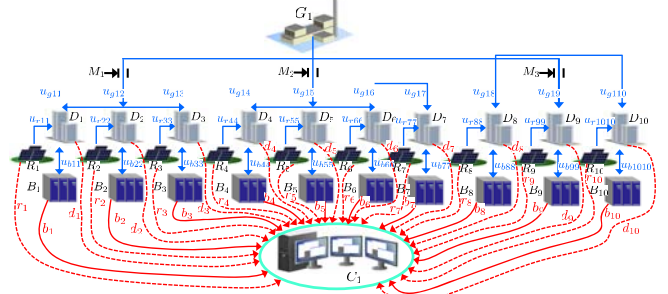
17



Distributed Control

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Simulation model — Predictive control, Central control



Need information on all variables in power network

→ When the network becomes large-scale, amount of calculation becomes huge.

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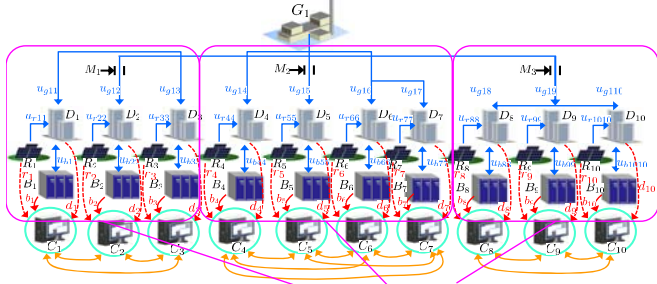
18



Distributed Control

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Block diagram — Distributed predictive control



There is no common weight variable
each calculator needs only local information

even if the network becomes large-scale,
amount of calculation of each calculator don't becomes huge.

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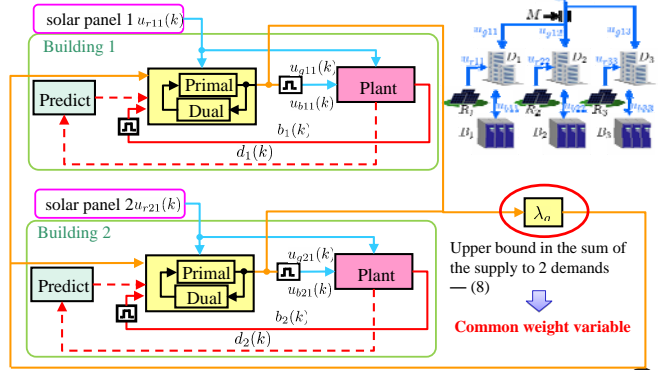
19



Distributed Control - Simulation block

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Block diagram — Predictive control



Upper bound in the sum of
the supply to 2 demands
— (8)

Common weight variable

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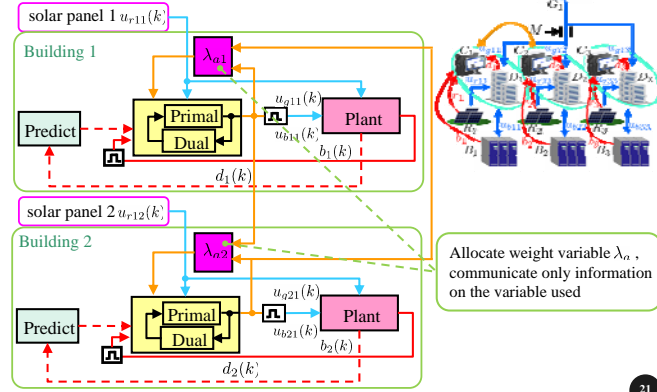
20



Distributed Control - Simulation block

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Block diagram — Distributed predictive control



Allocate weight variable λ_n .
communicate only information
on the variable used

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21



Distributed Control

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In power network, for a certain demand j

Utility function

$$(2) \rightarrow J_{b,i} = \sum_{k=0}^N \sum_{i \in \mathcal{N}_{BD,j}} w_{B,i} (b_i(k) - b_{i,ref})^2$$

$$(3) \rightarrow J_{g,j} = \sum_{k=0}^N \sum_{i \in \mathcal{N}_{GD,j}} w_{G,i} u_{gij}^2$$

$$(4) \rightarrow J_d = \sum_{k=0}^N \sum_{j \in \mathcal{D}} w_{D,j} \left(\sum_{i \in \mathcal{N}_{GD,i}} u_{gij}(k) + \sum_{i \in \mathcal{N}_{RD,i}} u_{rij}(k) + \sum_{i \in \mathcal{N}_{BD,i}} u_{bij}(k) - d_j(k) \right)^2$$

Constraints

$$(5) \rightarrow u_{gij}(k) \geq 0, u_{rij}(k) \geq 0$$

$$(6) \rightarrow \sum_{i \in \mathcal{N}_{BD,i}} u_{bij}(k) \geq - \sum_{i \in \mathcal{N}_{GD,i}} u_{gij}(k) - \sum_{i \in \mathcal{N}_{RD,i}} u_{rij}(k)$$

$$(7) \rightarrow r_i(k) = \sum_{j \in \mathcal{N}_{BD,i}} u_{rij}(k), i \in \mathcal{N}_{RD,j}$$

$$(8) \rightarrow \sum_{j \in \mathcal{D}, i} u_{gij} - x_i = M_i, x_i \geq 0$$

$$(9) \rightarrow b_i(k) \in [0, B_{i,max}], i \in \mathcal{N}_{BD,j}$$

coupling
communicate these information
to decide common variable

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22



Outline

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23



Simulation model

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Constant

$$\tau = 6[\text{min}]$$

$$q_b = 0.0005, q_d = 10, q_g = 0.1$$

$$B_{1,max} = 50[\text{kW}], b_{1,ref} = B_{1,max}/2$$

$$b_1(0) = 10[\text{kW}]$$

$$M_1 = 6.0[\text{kW}], M_2 = 9.0[\text{kW}], M_3 = 10.0[\text{kW}]$$

Utility function

use (2), (3), (4)

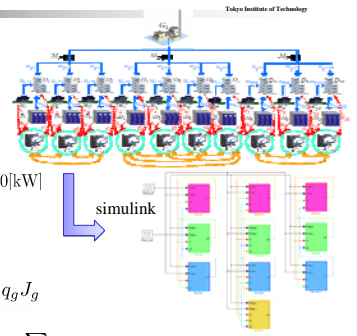
$$\min_{u(k), k=0, \dots, N} J = q_d J_d + q_b J_b + q_g J_g$$

Constraints

$$u_{gij}(k) \geq 0, u_{rij}(k) \geq 0 \quad \sum_{j \in \mathcal{D}, i} u_{gij} + x_i = M_i, x_i \geq 0$$

$$u_{bij}(k) \geq u_{gij}(k) - u_{rij}(k) \quad r_i(k) = \sum_{j \in \mathcal{N}_{BD,i}} u_{rij}, i \in \mathcal{N}_{RD,j}$$

$$b_i(k) \in [0, B_{i,max}] \quad j \in \mathcal{D}, k = 0, \dots, N$$

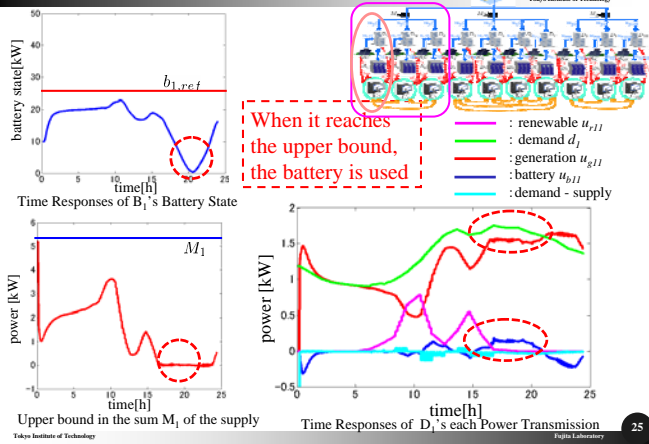


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24



Simulation



25



Conclusion

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- Introduction of generalization power network
- Introduction of distributed predictive control
- Application of distributed predictive control for a dynamic OPF
- validation of the utility of distributed predictive control by simulation

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26