

Discrete-time Lloyd Descent for Coverage Control



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- Algorithm
- Experiment
- Future work



Algorithm

○ Locational optimization

Let Q be convex polytope in \mathbb{R}^N ϕ : density function $P = \{p_1, p_2, \dots, p_n\}$: location of n sensors $W = \{W_1, W_2, \dots, W_n\}$: the partition of Q $f(\|q - p_i\|)$: distance function

consider the following objective function

$$H(P, W) = \sum_{i=1}^n \int_{W_i} f(\|q - p_i\|) \phi(q) dq \quad (1)$$



Minimizing the objective function

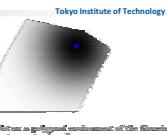


Fig. 1. Coverage problem of the sensor network generated by the Lloyd's algorithm.



Algorithm

○ Voronoi Partitions

Voronoi Partitions generated by the points $\{p_1, p_2, \dots, p_n\}$

$$V(P) = \{q \in Q \mid \|q - p_i\| \leq \|q - p_j\|, \forall j \neq i\} \quad (2)$$

Then, define

$$H_V(P) = H(P, V(P)) \quad (3)$$

Define the mass, centroid and polar moment of inertia about a region $V \subset \mathbb{R}^N$

$$M_{V'} = \int_{V'} \rho(q) dq \quad C_{V'} = \frac{1}{M_{V'}} \int_{V'} q \rho(q) dq$$

$$J_{V'} = \int_{V'} \|q - p\|^2 \rho(q) dq$$



Algorithm

Let $T : Q^n \rightarrow Q^n$ be a mapping satisfying properties (a) and (b)(a) for all $i \in \{1, \dots, n\}$, $\|T_i(P) - C_{V_i(P)}\| \leq \|p_i - C_{V_i(P)}\|$ where T_i denotes the i th component of T (b) If P is not centroidal, then there exists a j such that

$$\|T_j(P) - C_{V_j(P)}\| < \|p_j - C_{V_j(P)}\|$$

Then the sequence $\{T^m(P_0)\}_{m \in \mathbb{N}}$, if this set is finite, converge to a centroidal Voronoi configuration

Algorithm

Proof: about $H(P, W) = \sum_{i=1}^n \int_{W_i} f(\|q - p_i\|) \phi(q) dq$

using the parallel axis theorem

$$H(P, W) = \sum_{i=1}^n J_{W_i, C_{W_i}} + \sum_{i=1}^n M_{W_i} \|p_i - C_{W_i}\|^2 \quad (4)$$

Independent of P Dependent on P

therefore

$$H(P', W) \leq H(P, W) \quad (5)$$

as long as $\|p_i - C_{W_i}\| \leq \|p_i - C_{V_i}\|$ for all $i \in \{1, 2, \dots, n\}$



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Moreover, since Voronoi partition is the optimal one

$$H(P, V(P)) \leq H(P, W) \quad (6)$$

Because of property (a) of T

$$\underline{H(T(P), V(P))} \leq H(P, V(P)) = H_V(P) \quad (7)$$

And because of (6)

$$H_V(T(P)) = H(T(P), V(T(P))) \leq \underline{H(T(P), V(P))} \quad (8)$$

Hence

$$H_V(T(P)) \leq H_V(P) \quad (9)$$



H_V is a descent function for the algorithm T

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Experiment

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Dimension : 1D The number of agents: 3

$$p(k+1) = Ap(k) + B \begin{bmatrix} p_1 \\ p_2 \\ p_3 \end{bmatrix} = \frac{1}{4} \begin{bmatrix} 1 & 1 & 0 \\ 1 & 2 & 1 \\ 0 & 1 & 1 \end{bmatrix} \begin{bmatrix} p_1 \\ p_2 \\ p_3 \end{bmatrix} + \begin{bmatrix} 0 \\ 0 \\ 1/2 \end{bmatrix}$$

- 1:Update Voronoi partition and all $moving[i] = 1$
- 2:agent i move toward $p_i(k+1)$
- 3:if agent i arrive at $p_i(k+1)$ $moving[i] = 0$
- 4:if all $moving[i] = 0$, return to (1) until convergence

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Future Works

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- carry out the coverage control (1-D) experiment
- consider the result of the expriment

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